

Estimating Entropy and Divergence of Sensor Data

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Problem and Motivation

- Many simple, myopic sensors
- Would like to fuse myopic information to attain more global view of battlefield scenario
- Need to understand relationships between sensor outputs



Difficulties and Complications

- Need fast, low-complexity algorithms (but can exploit temporal information)
- Unknown scene and sensor geometry
- Complex, dynamic environment
- Multiple, widely separated, possibly dynamic sensors
- Uncalibrated, possibly multi-modal, sensors (unknown parameters, geometry)
- Noise



A Key Task and Possible Approach

- Correspondence and fusion are
 - difficult with many parameters (or nonparametric) and with nonlinear transformations
 - meaningless without common information in data streams
- A key task is to determine how related are the outputs of two sensors
- A possible approach:
 - Determine sensor pairs/groups that are highly related
 - Perform correspondence and fusion first with maximally related sensors
 - Incorporate other sensor outputs



A "Distilled" Problem

- The Problem: How to estimate the entropy and divergence of two sources based only on one realization from each source ?
- Assumption: Both are finite-alphabet, finitememory, stationary sources.
- Our goal: Want good estimates, fast convergence, and reasonable computational complexity.



Definitions

Entropy and Divergence rates are defined as follows:

$$H(q_z) = \lim_{n \to \infty} \frac{1}{n} E[\log \frac{1}{q_z(z^n)}]$$

$$= \sum_{s \in S} q_z(s) \sum_{a \in \chi} q_z(a \mid s) \log \frac{1}{q_z(a \mid s)}.$$

$$D(q_z \parallel p_x) = \lim_{n \to \infty} \frac{1}{n} E[\log \frac{q_z(z^n)}{p_x(z^n)}]$$

$$= \sum_{s \in S} q_z(s) \sum_{a \in \chi} q_z(a \mid s) \log \frac{q_z(a \mid s)}{p_x(a \mid s)}.$$

Possible Approaches and Previous Work

- Model Based: Assume model to get empirical distribution. Then plug-in to formulas.
- Universal Algorithms: Use universal compression algorithm or related methods to estimate entropy. E.g., can use Lempel-Ziv string matching method inspired by LZ data compression algorithm (Wyner, Ziv '89, Kontoyiannis et al. '94,'96,'98, Quas '95, Shields '92).

 $L_n = 1 + \text{length of the longest prefix that recurs}$

$$\hat{H_n} = \frac{\log_2 n}{L_n}$$

e.g., for "abaabaaa", $L_8=5$.

 Very little previous work on divergence estimation (Ziv, Merhav '93)



Weaknesses and Overcoming Them

- Model-based methods need to know order of model and are not universal.
- Existing universal methods (e.g., LZ-based methods) converge very slowly.
- Can we retain advantages of both?
- Recent interest in Burrows-Wheeler Transform (block sorting) for data compression (Burrows, Wheeler '94)
- Can this be adapted for entropy and divergence estimation?

Burrows-Wheeler Block Sorting

- 1. Get every shift of the input sequence.
- 2. Sort them alphabetically.
- 3. Output the last column of the sorted table, as well as the position of the original sequence.
- Reversible transform.
- Can be implemented in *O(n)* time/space complexity.
- Sorts input sequence according to context.
- Output is close to a piecewise i.i.d. distribution (Effros, Visweswariah, Kulkarni, Verdu '02).



BWT: Example

	EOF symbol	sorted table	output:
input →	banana\$	anana\$b	b
cyclic shifts	\$banana	ana\$ba n	n
	a\$banan sort	a\$bana n	output N
	na\$bana \Longrightarrow	banana\$	⇒\$
	ana\$ban	nana\$b a	а
	nana\$ba	na\$ban a	а
ĺ	anana\$b	\$banana	а



Entropy Estimation – Basic Idea

The BWT sorts the sequence according to context, and the output is close to piecewise i.i.d.





Entropy Estimation via BWT

- 1. Reverse original sequence \mathbf{z}^n .
- 2. Run BWT on reversed sequence.
- Divide BWT output into segments, according to transitions of distribution.
- 4. Estimate letter probabilities in each i.i.d. segment.
- 5. Combine estimates in each segment (average log of probability of each segment) to calculate probability of zⁿ and in turn the entropy.



Entropy Estimation via BWT: cont.



Block diagram of the entropy estimator



Divergence Estimation – Basic idea

Recall,

$$D(q \parallel p) = \lim_{n \to \infty} \frac{1}{n} \left(\log \frac{1}{p_x(z^n)} - \log \frac{1}{q_z(z^n)} \right)$$

- Second term is just entropy of **z**.
- First term (cross term) is probability of z, but according to distribution for x.
- Key idea to estimate cross term is *joint* BWT of both x and z.



Estimation of Cross Term

- Joint BWT of **x** and **z**:
 - 1. Concatenate **x** and **z**, adding '\$' to the end of each.
 - Sort the table of cyclic shifts. Note whether each letter comes from x or from z during the BWT sorting.
 - 3. Output the last column.

This gather letters with same context from both **x** and **z** together.

- Segment the joint BWT output *according to* **x**.
- Compute probability of z using probabilities estimated according to x.



Estimation of Cross Term – continued





Divergence Estimation via BWT

- 1. Estimate entropy $H_z = -1/n \log q_z(z^n)$.
- Estimate the cross term $-1/n \log p_x(z^n)$.
- 3. Subtract the entropy term from the cross term.



Block diagram of the divergence estimator



Approaches to Segmentation

- Uniform segmentation: divide the BWT output sequence into equal-length segments.
- Adaptive segmentation: make a new segment only when we detect a transition. The number and length of segments are adapted to the source.



Uniform Segmentation

Divide the sequence (of length n) into equal-length segments of length w(n). Want those segments containing transitions to be negligible, but don't want too many segments.





Adaptive Segmentation

- Uniform segmentation ends up with many more segments than the actual number of states.
- Instead of using equal-length segments, estimate positions of transitions, and make new segment only when we detect a transition.
- Two-level blocks are introduced. In level-1 (with block length k₁), roughly locate the positions. In level-0 (with block length k₀), refine our estimate of the positions.



Experimental Results

- Have tested algorithm on simulated date (randomly generated binary tree sources).
- Compared new algorithm with LZ-based methods for both entropy and divergence estimation.
- Compared new algorithm with empirical distribution plug-in scheme for both entropy and divergence.
- Work with text files in progress, and sensor data forthcoming.



New algorithm vs. LZ – Entropy Estimator





New algorithm vs. LZ – Divergence Estimator





New algorithm vs. Plug-in – Entropy Estimator





New algorithm vs. Plug-in – Divergence Estimator





Experimental Results

- New algorithm converges much faster than LZ-based algorithm.
- Choosing "right order" is critical for the empirical distribution plug-in scheme.
- New algorithm has an intrinsic advantage by not assuming $|S| = |\chi|^{D}$. Also, doesn't assume prior knowledge about memory length or number of states.

Preliminary Thoughts on Estimating Mutual Information

- Mutual information tells how much one information one source provides about another.
- Can estimate mutual information via entropy or divergence:
 - I(X;Y) = H(X) H(X|Y) = H(X) + H(Y) H(X,Y).
 - I(X;Y) = D(p(x,y) || p(x)*p(y)).
- Alphabet size increases complexity.



Summary

- A new entropy and divergence estimator based on BWT (block sorting).
- Doesn't require knowledge of distribution or parameters of the sources.
- Efficient algorithm, good estimates, fast convergence.
- Significantly outperforms other algorithms tested.
- Expect this to be useful in a wide range of applications --- specifically, a key component in general correspondence and fusion algorithms.



Future Work

- Fine-tune and improve algorithm (e.g., segmentation procedure).
- Further analysis of performance.
- Extensions (e.g., continuous source, nonstationarity).
- Assess performance on actual sensor data.
- Implement algorithms for mutual information.
- Integrate as a component of general correspondence and fusion algorithms.